

# A live data synchronization using clsync

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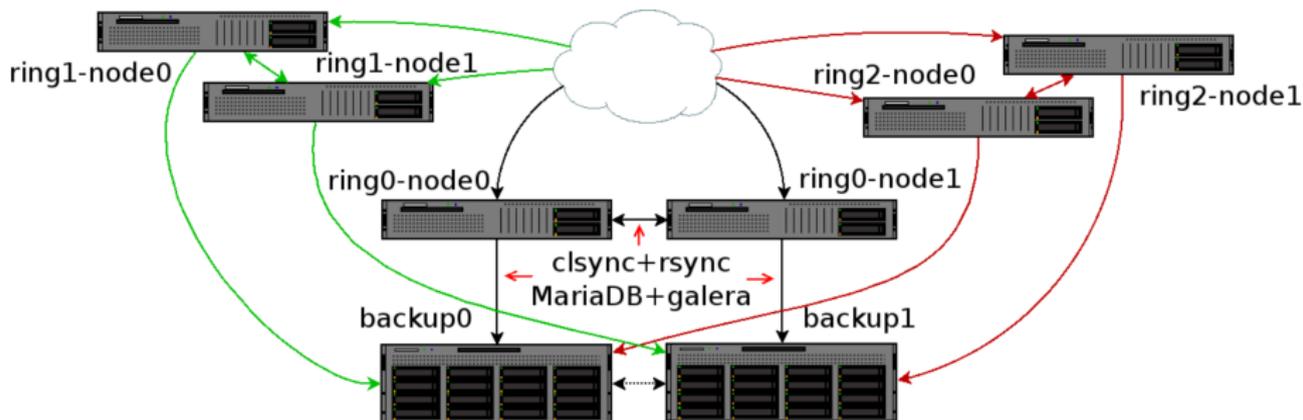


Our task:

- Deploy HA/LB clusters:
  - hosting
  - VoIP
  - corporate services
- HPC systems management and synchronization
- Backups of all the above
- ... and we have just limited hardware resources



# Infrastructure



- Separation by rings based on importance and trust
- HA and backups for each ring



# Synchronization approaches

## Possible solutions:

- RO file systems
  - limited application area
- block replication
  - unacceptable performance (DRBD + OCFS2)
  - not resistant to split brain
- network file systems (e.g. CEPH)
  - high latency (on 1 Gb/s)
  - kernel panics (year 2012)
- file level replication

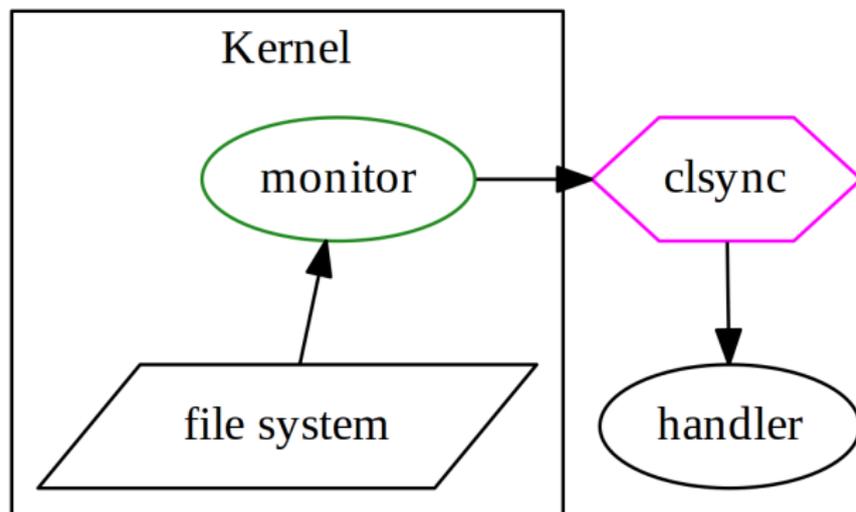


# Choosing software

- Isyncd: best match, but:
  - high CPU load ( $> \frac{1}{2}$  code on LUA)
  - bugs and complexity of LUA code support
  - no threading, bad for large dirs ( $> 10^6$  objects)
  - limited events aggregation
  - no \*.so support
  - no BSD support
- incron
  - no recursion, no events
- csync2
  - no events, extremely slow
- librnotify
  - not existed at that time :)
  - bare inotify interface



# How clsync works



Handler can be anything, but common case is rsync.



# Choosing monitoring subsystem

Linux:

## dnotify

- + can trace any event
  - individual files can't be monitored
  - umount is blocked
  - signal-based notification
  - requires `man stat()`, only `fd` is provided

## inotify [choice]

- + epoll-based notification
- + full event info provided
  - no recursion
  - watch  $\Rightarrow$  path map required

# Choosing monitoring subsystem

## fanotify

- + recursion support
- + returns fd and pid *Rightarrow* path known
- no support for delete, move, rename events ;-(

## FreeBSD

- kqueue/kevent
- libinotify (using kqueue)
- BSM API (not original purpose)
- dtrace (unusable, path can't be extracted)

There are no worthy Linux inotify replacements in FreeBSD.  
Best match is kqueue interface.

# How clsync works

On default settings:

- 1 Initialization
- 2 Set inotify watches
- 3 Full file tree sync
- 4 Resync for new events
- 5 Wait, event aggregation and ↑



synchandler modes:

- *simple* — per event per file app call
- *direct* — per event app call
- *shell* — shell call on any sync event
- *rsyncdirect* — direct rsync call
- *rsyncshell* — call for rsync wrapper
- *rsyncso* — `clsyncapi_rsync()` callback with listing suitable for rsync
- *so* — load \*.so with generic `clsyncapi_sync()` callback using simple file listing

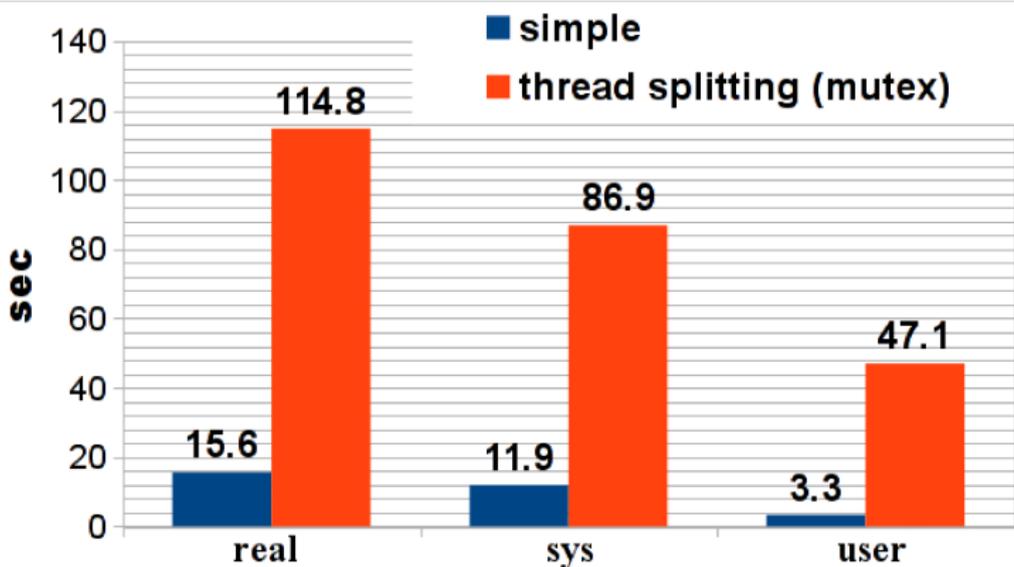


Security features optionally engaged:

- privilege drop
- use of capabilities
- namespace isolation
- cgroups isolation
- thread splitting (priv + common)
- seccomp isolation
  - mprotect() banned  $\Rightarrow$  no thread splitting
- process splitting (priv + common)



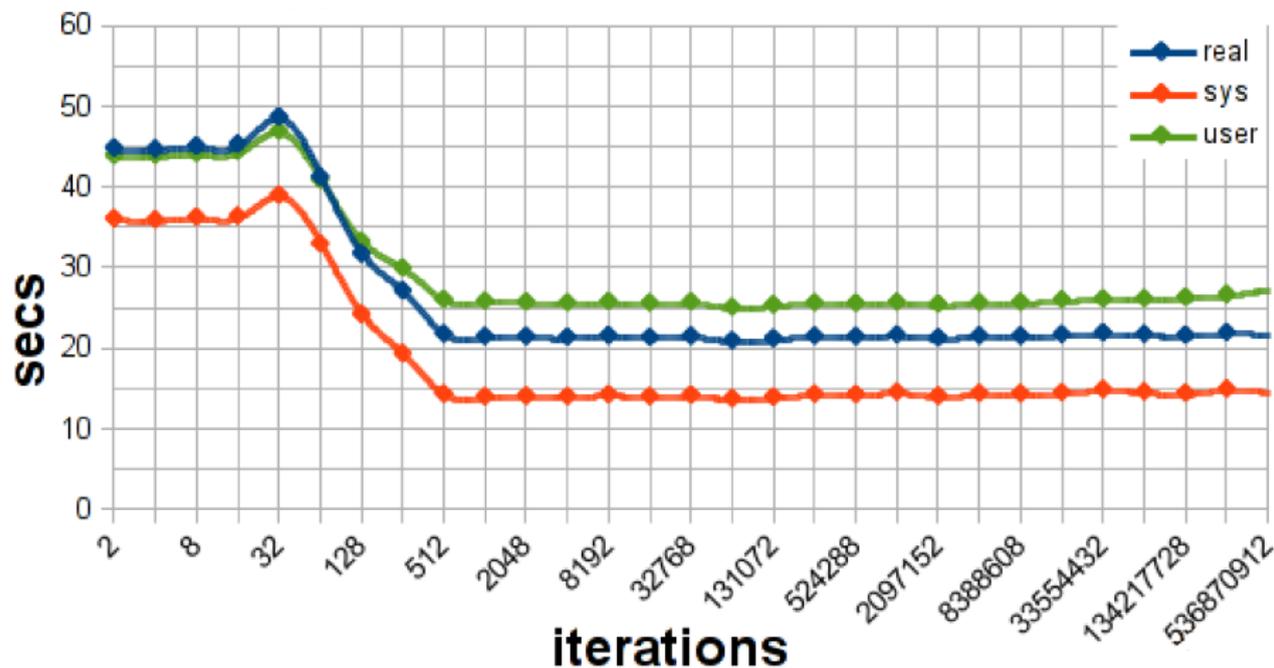
# Threading performance



pthread mutex is not designed for high-speed locks



# Spinlock to mutex switch

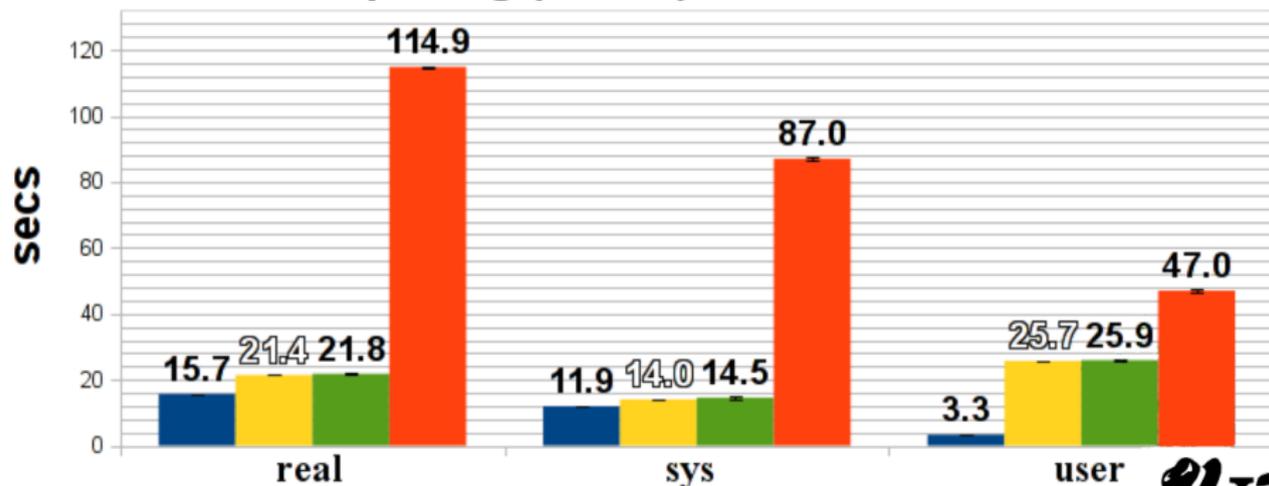


tested on init sync procedure



# Smart locks

- simple
- thread splitting (high load locks + auto adjust)
- thread splitting (high load locks)
- thread splitting (mutex)

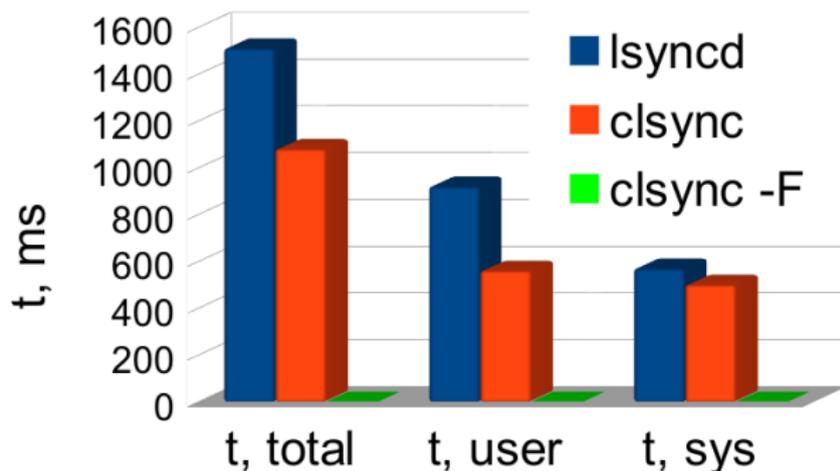


# More features

- threading support (thread per sync event)
- regexp support and fs object type selection
- UNIX-socket control interface
- auto switch between spin\_lock и mutex
- fast initial sync



# Overhead measurement



- lsyncd v.2.1.x
- clsync v.0.4.x
- 4789558 files and dirs on tmpfs
- -F allows to bypass rules on initial sync



## Mirror directory

```
clsync -Mrsyncdirect -W/path/to/source_dir \  
-D/path/to/destination_dir
```

## One-time sync

```
clsync --exit-on-no-events --max-iterations=20 \  
--mode=rsyncdirect -W/var/www_new -Srsync -- \  
%RSYNC-ARGS% /var/www_new/ /var/www/
```



## Live correction of web site perms

```
clsync -w1 -t1 -T1 -x1 \  
-W/var/www/site.example.org/root \  
-Mdirect -Schown --uid 0 --gid 0 -Ysyslog -b1 \  
--modification-signature uid,gid -- \  
--from=root www-data:www-data %INCLUDE-LIST%
```

## Troubleshooting

```
Cannot inotify_add_watch() on [...]:  
No space left on device (errno: 28)
```

Increase sysctl **fs.inotify.max\_user\_watches**:

- one watch per dir
- no recursion
- kernel default is 8192

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# Status

- Main functionality is implemented
- Maintenance mode
- Latest release 0.4.2 a week ago ; -)

Distro support:

|          |                        |
|----------|------------------------|
| Gentoo   | official (w/o *-fbbsd) |
| Debian   | official (w/o fbbsd)   |
| RH-based | RPM available          |
| FreeBSD  | ports available        |

The only mandatory dep is glib.



# Summary

## Contacts:

|        |   |
|--------|---|
| Github | <a href="https://github.com/xaionaro/clsync">https://github.com/xaionaro/clsync</a> |
| IRC    | Freenode: #clsync   |
| e-mail | dyokunev@ut.mephi.ru<br>aasavchenko@ut.mephi.ru                                     |

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- Dmitry Yu. Okunev — main clsync author
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- oldlaptop and Enrique Martinez for their help

Thank you for your attention!



Backup slides follow...



## What we have:

- Limited hardware:
  - About 10 blades
  - 1 Gb/s interconnect (not very stable)
  - short on RAM and HDDs
- > 100 tasks (containers):
  - main university web site
  - departments sites
  - VoIP
  - VPN and Wi-Fi for students
  - internal services (e-mail, ntp, sks,...)



What can be done?

- Virtual machines  $\Rightarrow$  containers (LXC)
- Disk data deduplication:
  - base image + aufs/overlays
- Cheap HA/LB clusters

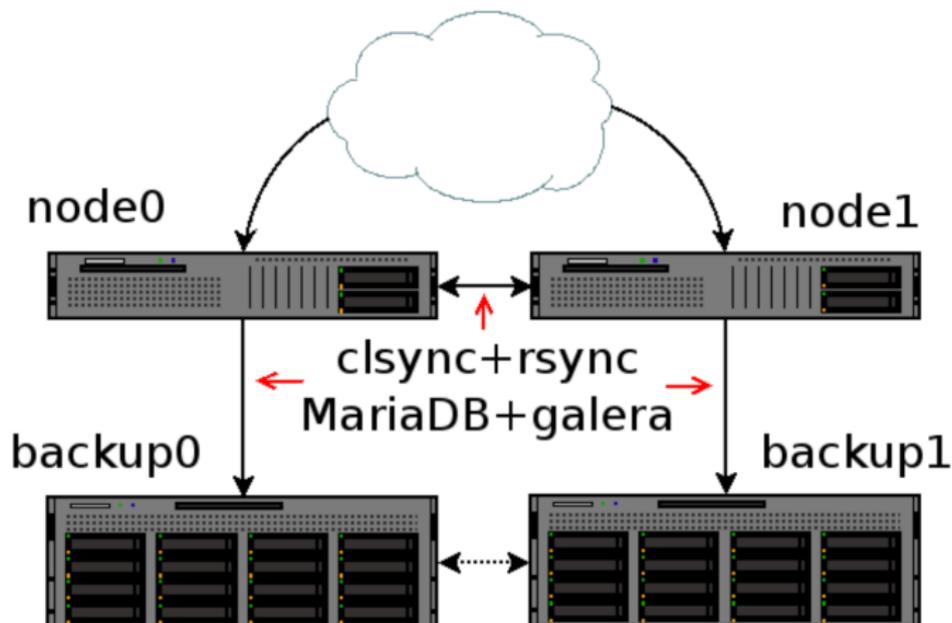


## Requirements:

- performance (minimal overhead)
- high availability (max few seconds downtime)
- reliability (failure minimization)
- versatility (wide range of use cases)
- fine tuning over aggregation



# Infrastructure core



- LXC
- Ext4 c clsync + rsync
- Percona



# Spinlock on single core

